

# **Quantum Information: Part I**

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# Quantum Information

- What is quantum information?
- Reply: Quantum information is information represented in quantum systems.
- What is information?
- What is a quantum system?

# Information

- Information is difference of entropies
- What is entropy?

# Entropy



$$S = k \cdot \log W,$$

where  $k$  is a constant,  $W$  is the number of microstates corresponding to a macroscopic state

# Entropy

“Elementary” entropy  $H(n)$  = number of elementary units (bits, trits, etc.) to coin  $n$  (uniform) conditions.

- Binary entropy:  $\{1, 2\} \mapsto \{0, 1\}$ ,
- $\{1, 2, 3\} \mapsto \{0, 1, 00\}$ ,
- $\{1, 2, 3, 4\} \mapsto \{00, 01, 10, 11\}$ , etc.
- $H_2(n) = \log_2 n = \frac{1}{\log 2} \log n$
- $H_3(n) = \frac{1}{\log 3} \log n$ , etc.
- Measure of “uncertainty”

# Entropy

Boltzmann: Identical particles with same internal condition indistinguishable.

- Let  $l$  be the number of particles, each having  $n$  potential conditions  $\{1, 2, \dots, n\}$ ,  $l \gg n$ .
- List the conditions of all particles:  $c_1 c_2 \dots c_l$ ,  
 $c_i \in \{1, 2, \dots, n\}$
- Assume condition  $i$  occurs  $k_i$  times, so  $k_1 + \dots + k_n = l$  and  $p_i = \frac{k_i}{l}$  is the probability (frequency) of condition  $i$

# Entropy

There are

$$\frac{l!}{k_1! \dots k_n!}$$

such lists (strings of conditions)

Entropy:

$$K \log \frac{l!}{k_1! \dots k_n!}$$

Per particle:

$$\frac{K}{l} \log \frac{l!}{k_1! \dots k_n!}$$

# Entropy

Stirling:  $\log k! = k \log k - k + O(\log k)$ , so

$$\begin{aligned} & \frac{K}{l} \log \frac{l!}{k_1! \dots k_n!} \\ = & \frac{K}{l} (l \log l - l + O(\log l)) \\ & - \sum_{i=1}^n (k_i \log k_i - k_i + O(\log k_i)) \\ = & -K \sum_{i=1}^n p_i \log p_i + O\left(\frac{\log l}{l}\right). \end{aligned}$$

# Entropy

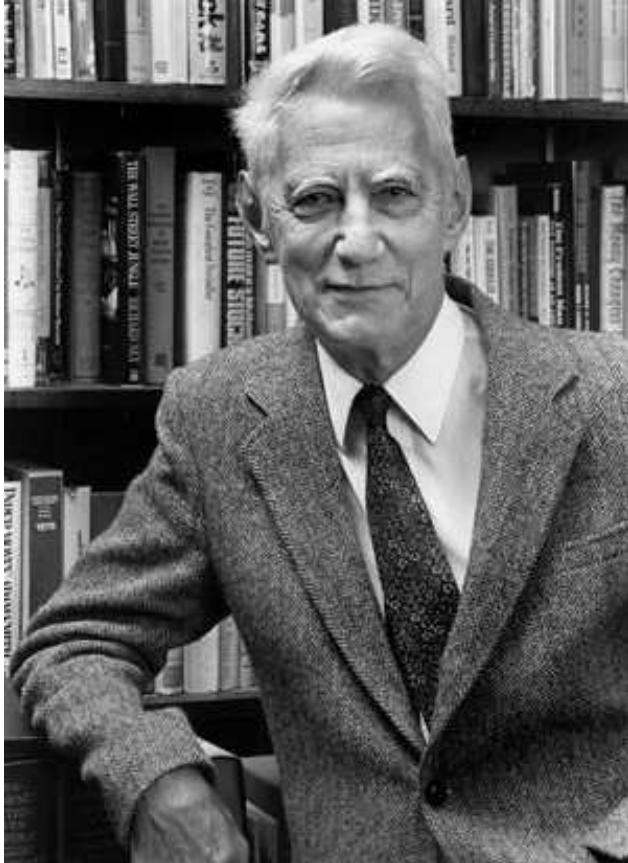
For a probability distribution  $(p_1, \dots, p_n)$  of events  $\{e_1, \dots, e_n\}$ , define

$$H(p_1, \dots, p_n) = -K \sum_{i=1}^n p_i \log p_i$$

For  $p_1 = \dots = p_n = \frac{1}{n}$

$$H\left(\frac{1}{n}, \dots, \frac{1}{n}\right) = -K \cdot n \cdot \frac{1}{n} \log \frac{1}{n} = K \log n$$

# Entropy



- $H(p_1, \dots, p_n)$  symmetric, continuous
- $H(\frac{1}{n}, \dots, \frac{1}{n})$  non-negative, strictly increasing in  $n$
- $H(p_1, \dots, p_n) + p_n H(q_1, \dots, q_m) = H(p_1, \dots, p_{n-1}, p_n q_1, \dots, p_n q_m)$

$$\Rightarrow H(p_1, \dots, p_m) = -K \sum_{i=1}^n p_i \log p_i$$

Claude Shannon (1916-2001)

# Joint Entropy

Let

$$X = \{x_1, \dots, x_n\}$$

be a random variable with distribution  $p(x_1), \dots, p(x_n)$ . Then

$$H(X) = - \sum_{i=1}^n p(x_i) \log p(x_i).$$

If also  $Y = \{y_1, \dots, y_m\}$  is a random variable, the *joint entropy* is

$$H(X, Y) = - \sum_{i=1}^n \sum_{j=1}^m p(x_i, y_j) \log p(x_i, y_j)$$

# Conditional Entropy

$$H(X, Y) = - \sum_{i=1}^n \sum_{j=1}^m p(x_i, y_j) \log p(x_i, y_j),$$

but if  $Y$  is known to assume value  $y_j$ , then

$$H(X \mid y_j) = - \sum_{i=1}^n p(x_i \mid y_j) \log p(x_i \mid y_j),$$

and

$$H(X \mid Y) = \sum_{j=1}^m p(y_j) H(X \mid y_j).$$

“Uncertainty of  $X$  when  $Y$  is known”

# Conditional Entropy

Lemma:  $H(X | Y) \leq H(X)$

Lemma:  $H(X, Y) \leq H(X) + H(Y)$

Lemma:  $H(X | Y) = H(X, Y) - H(Y)$

(Uncertainty of  $X$  when  $Y$  is known)

# Conditional Entropy

Example: Team  $A$  wins with probability  $\frac{1}{2}$ ,  $X = \{\text{win}, \text{loss}\}$ .

Then  $H(X) = -\left(\frac{1}{2} \log_2 \frac{1}{2} + \frac{1}{2} \log_2 \frac{1}{2}\right) = 1$ .

As a *home team*,  $A$  wins with  $\frac{3}{4}$  probability, but as *visitor*,  $A$  wins only with  $\frac{1}{3}$  probability.

$$H(X \mid h) = -\left(\frac{3}{4} \log_2 \frac{3}{4} + \frac{1}{4} \log_2 \frac{1}{4}\right) = 0.811278\dots,$$

$$H(X \mid v) = -\left(\frac{1}{3} \log_2 \frac{1}{3} + \frac{2}{3} \log_2 \frac{2}{3}\right) = 0.918296\dots$$

# Conditional Entropy

$$H(X \mid h) = -\left(\frac{3}{4} \log_2 \frac{3}{4} + \frac{1}{4} \log_2 \frac{1}{4}\right) = 0.811278\dots,$$

$$H(X \mid v) = -\left(\frac{1}{2} \log_2 \frac{1}{2} + \frac{1}{2} \log_2 \frac{1}{2}\right) = 0.918296\dots$$

Let  $Y = \{0, 1\}$  be a fair coin toss for deciding if team  $A$  plays home. Then

$$H(X \mid Y) = \frac{1}{2}H(X \mid h) + \frac{1}{2}H(X \mid v) = 0.864787\dots$$

# Information

Mutual information of  $X$  and  $Y$ :

$$I(X : Y) = H(X) - H(X \mid Y)$$

$$\begin{aligned} I(X : Y) &= H(X) - H(X \mid Y) \\ &= H(X) - (H(X, Y) - H(Y)) \\ &= H(X) + H(Y) - H(X, Y) \\ &= I(Y : X) \end{aligned}$$

“Uncertainty of  $X$  minus uncertainty of  $X$  when  $Y$  known”

Previous example:

$$I(X : Y) = 1 - 0.864787 \dots = 0.135213 \dots$$

# Quantum Information



- Quantum entropy by Gedanken Experiment (1927)
- Coincides with Shannon (and Boltzmann) entropy on classical systems

John von Neumann (1903–1957)

# Mechanics

Newtonian equation of motion:

$$F = ma = m \frac{d}{dt} v = \frac{d}{dt} mv = \frac{d}{dt} p$$

Total energy:

$$H = \frac{1}{2}mv^2 + V(x) = \frac{p^2}{2m} - \int_{x_0}^x F(s) ds$$

Hamiltonian reformulation:

$$\frac{d}{dt}x = \frac{\partial}{\partial p} H, \quad \frac{d}{dt}p = -\frac{\partial}{\partial x} H$$

# Mechanics

Classical:

$$\frac{d}{dt}x = \frac{\partial}{\partial p}H, \quad \frac{d}{dt}p = -\frac{\partial}{\partial x}H$$

Quantum:

$$\frac{\partial}{\partial t}\psi = -iH\psi,$$

where  $\psi$  is the *wave function*.

# Wave Function

Max Born's interpretation:

$$|\psi(x, t)|^2$$

is the probability density of the particle position at time  $t$ :

$$\mathbb{P}(a \leq x \leq b) = \int_a^b |\psi(x, t)|^2 dx$$

On the other hand (omitting  $t$ ):

$$\widehat{\psi}(p) = \mathcal{F}[\psi(x)](p) = \int_{-\infty}^{\infty} \psi(x) e^{-2\pi i x p} dx$$

is the probability density of the particle momentum.

# Wave Function

On the other hand (omitting  $t$ ):

$$\widehat{\psi}(p) = \mathcal{F}[\psi(x)](p) = \int_{-\infty}^{\infty} \psi(x) e^{-2\pi i x p} dx$$

is the probability density of the particle momentum:

$$\mathbb{P}(a \leq p \leq b) = \int_a^b |\widehat{\psi}(p)|^2 dp$$

$\psi$  gives the complete characterization of the system at a fixed time

# Finite Quantum Systems

- Nuclear spin
- Photon polarization

Wavefunction  $\psi$  defined on a finite set.

Formally,

$$\psi = \alpha_1 \psi_1 + \alpha_2 \psi_2 + \dots + \alpha_n \psi_n,$$

where  $\{\psi_1, \dots, \psi_n\}$  is an orthonormal basis of  $n$ -dimensional complex vector space.

For *mixed states*, representation must be generalized.

# Formalism of Quantum Mechanics



- Hilbert space
- Linear mappings (operators)

John von Neumann (1903–1957)

# Formalism



Bra-ket notions

$$\langle x | y \rangle, |y\rangle, \langle x|, |y\rangle\langle x|,$$

Paul Dirac (1902-1984)

# Formalism

$n$ -level system  $\leftrightarrow n$  perfectly distinguishable values

Formalism based on  $H_n \simeq \mathbb{C}^n$  ( $n$ -dimensional Hilbert space)

- Hermitian inner product  $\langle x | y \rangle = x_1^* y_1 + \dots + x_n^* y_n$
- Norm  $\|x\| = \sqrt{\langle x | x \rangle}$
- Ket-vector  $|x\rangle = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}$
- Bra-vector  $\langle x| = (|x\rangle)^* = (x_1^*, \dots, x_n^*)$
- Adjoint matrix:  $(A^*)_{ij} = A_{ji}^*$  for  $m \times n$  matrix  $A$

# Formalism

- Trace:  $\text{Tr}(A) = \sum_{i=1}^n A_{ii}$
- For orthonormal basis  $\{x_1, \dots, x_n\}$ ,  
$$\text{Tr}(A) = \sum_{i=1}^n \langle x_i | Ax_i \rangle$$
- Positivity:  $A \geq 0$  iff  $(\forall x) \langle x | Ax \rangle \geq 0$
- Self-adjointness:  $A^* = A$
- Unitarity:  $UU^* = U^*U = I$
- Normality:  $A^*A = AA^*$

# Formalism

Kronecker product:

$$A = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1s} \\ a_{21} & a_{22} & \dots & a_{2s} \\ \vdots & \vdots & \ddots & \vdots \\ a_{r1} & a_{r2} & \dots & a_{rs} \end{pmatrix}, B = \begin{pmatrix} b_{11} & b_{12} & \dots & b_{1u} \\ b_{21} & b_{22} & \dots & b_{2u} \\ \vdots & \vdots & \ddots & \vdots \\ b_{t1} & b_{t2} & \dots & b_{tu} \end{pmatrix}$$

$$A \otimes B = \begin{pmatrix} a_{11}B & a_{12}B & \dots & a_{1s}B \\ a_{21}B & a_{22}B & \dots & a_{2s}B \\ \vdots & \vdots & \ddots & \vdots \\ a_{r1}B & a_{r2}B & \dots & a_{rs}B \end{pmatrix}$$

# Formalism

- $|x\rangle\langle y| = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \otimes (y_1^*, \dots, y_n) =$   
$$\begin{pmatrix} x_1 y_1^* & x_1 y_2^* & \dots & x_1 y_n^* \\ \vdots & \vdots & \ddots & \vdots \\ x_n y_1^* & x_n y_2^* & \dots & x_n y_n^* \end{pmatrix}$$
- $|x\rangle\langle y| |z\rangle = \langle y | z \rangle |x\rangle$
- If especially  $\|x\| = 1$ ,  $|x\rangle\langle x|$  is a projection onto a subspace generated by  $x$ .

# Formalism

- Each normal  $A$  has spectral representation

$$A = \lambda_1 |x_1\rangle\langle x_1| + \dots + \lambda_n |x_n\rangle\langle x_n|,$$

where  $\{x_1, \dots, x_n\}$  is an orthonormal basis of  $H_n$  and  $\lambda_1, \dots, \lambda_n$  the eigenvalues of  $A$ .

- If  $A$  is self-adjoint, each  $\lambda_i \in \mathbb{R}$
- If  $A$  is unitary, each  $\lambda_i$  has  $|\lambda_i| = 1$
- If  $A$  is positive, each  $\lambda_i \geq 0$ .
- $\text{Tr}(A) = \lambda_1 + \dots + \lambda_n$ .

# Structure of Quantum Mechanics

- State of a physical system: Unit-trace, positive operator  $T$ :

$$T = \lambda_1 |\mathbf{x}_1\rangle\langle\mathbf{x}_1| + \dots + \lambda_n |\mathbf{x}_n\rangle\langle\mathbf{x}_n|,$$

where  $\lambda_i \geq 0$ ,  $\lambda_1 + \dots + \lambda_n = 1$  (density matrix).

- Observable: Self-adjoint operator  $A$ :

$$A = \mu_1 |\mathbf{y}_1\rangle\langle\mathbf{y}_1| + \dots + \mu_n |\mathbf{y}_n\rangle\langle\mathbf{y}_n|,$$

where  $\mu_i \in \mathbb{R}$  are the potential values of  $A$

- Minimal interpretation:

$$\mathbb{P}(\mu_i) = \text{Tr}(T |\mathbf{y}_i\rangle\langle\mathbf{y}_i|)$$

is the probability of seeing value  $\mu_i$  if  $A$  is observed when the system is in state  $T$ .