

A NOTE ON SQUARE-FREE SHUFFLES OF WORDS

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Shuffles and repetitions

- Both popular items in words and languages.
- Prodinge, Urbanek (1979) considered *binary* words avoiding long squares using (perfect) shuffles.

These cases for binary words were continued by Rampersad (2004) and Currie, Rampersad (2010).

A **shuffle** of two words u and v :

$$u = u_1 u_2 \cdots u_n \quad \text{and} \quad v = v_1 v_2 \cdots v_n$$

with $u_i, v_i \in A^*$:

$$u_1 v_1 u_2 v_2 \cdots u_n v_n \in u \sqcup v .$$

It is a **perfect shuffle** if u_i and v_i are single letters:

$$aabb \sqcup_P bbab = ababbabb$$

Self-shuffle

- Charlier, Kamae, Puzynina, Zamboni (2013):
infinite word w obtained by shuffling it with itself:

$$w \in w \sqcup w.$$

- [CKPZ] A neat proof for:
the Fibonacci word f can be self-shuffled.

Indeed, $f = \varphi(f)$ for $\varphi(0) = 01$, $\varphi(1) = 0$.

Hence $\varphi^2(a) = \varphi(a)a$ for $a \in \{0, 1\}$. So

$$\begin{aligned} f &= a_1 a_2 \cdots = \varphi(a_1) \varphi(a_2) \cdots = \varphi^2(a_1) \varphi^2(a_2) \cdots \\ &= \varphi(a_1) a_1 \varphi(a_2) a_2 \cdots \in f \sqcup f. \end{aligned}$$

- [CKPZ]: *0f cannot be self-shuffled.*
- [CKPZ]: A longer proof for:
the Thue-Morse can be self-shuffled.

Theorem [CKPZ]. *For self-shuffled infinite $w \in w \sqcup w$, each sufficiently long prefix u is *Abelian bordered*:*

$$u = xvx' \quad (x, x' \text{ have equal Parikh vectors})$$

Now what?

Theorem. *There exist infinite square-free words u, v over a 4-letter alphabet such that the perfect shuffle $u \sqcup_P v$ is square-free.*

Of course, this is not self-shuffling.

In a 3-letter alphabet the above is not possible, **but:**

Theorem. *There exists an infinite square-free ternary word u such that $u \sqcup u$ contains a square-free word.*

Lengths: Problem solved

Theorem [TH, Mike Müller (2013) – ArXiv].

For each $n \geq 3$, there exists a square-free ternary word u of length n such that $u \sqcup u$ contains a square-free word.

Notation in the paper

Instead of \sqcup we used **conducting sequences** $\beta \in \{0, 1\}^*$ that pick letters from the two given words: if u and v are finite words of length n , then

$$u \sqcup v = \{ \beta[u, v] \mid \beta \in \{0, 1\}^{2n} \}.$$

For infinite words u and v , the binary β is also infinite. Also $\beta[u] = \beta[u, u]$.

Example. Let $u = 0100$, and $\beta = 00101110$. Then

$$\beta[u] = 01001000 \in u \sqcup v.$$

Perfect Shuffles

Example. Let $u, v \in \{0, 1, 2\}^*$ be square-free of length ≥ 3 . Then the perfect shuffle $u \sqcup_P v$ is *not square-free*.

Indeed, let without restriction

$$u = 01a_3 \cdots \quad \text{and} \quad v = b_1b_2b_3 \cdots .$$

Now $u \sqcup_P v = 0b_11b_2a_3b_3 \cdots$. To avoid squares, $b_1 = 2$ and $b_2 = 0$, and so $a_3 = 2$. Finally, $b_3 = 1$, and $u \sqcup_P v$ starts with the square 021021 ; a contradiction.

Four letters

Theorem. *There are square-free words $u, v \in \{0, 1, 2, 3\}^{\mathbb{N}}$ such that the perfect shuffle $u \sqcup_P v$ is square-free.*

Sketch. **Dean** (1965): Let u be an infinite square-free word **reduced** in the free group of two generators: u avoids the factors 02, 20, 13 and 31.

Let u' be the **dual** of u obtained by $0 \leftrightarrow 2$ and $1 \leftrightarrow 3$.

E.g., if $u = 0123$ then $u' = 2301$.

For any reduced square-free u , $u \sqcup_P u'$ is square-free.
(Details omitted.) □

Dean words

The following morphism (due to Thue) is square-free

$$\alpha(0) = 02102$$

$$\alpha(1) = 010212$$

$$\alpha(2) = 0121012$$

Thus so is the uniform $\sigma: \{0, 1, 2\}^* \rightarrow \{0, 1, 2, 3\}^*$

$$\sigma(0) = 03210323$$

$$\sigma(1) = 01032123$$

$$\sigma(2) = 01210123$$

Hence for each square-free $w \in \{0, 1, 2\}^*$, the image $\sigma(w)$ is square-free ‘Dean word’.

Single Ternary Words: $u \sqcup u$

Example. A square-free $w \in u \sqcup u$ may be obtained in several ways (using different conducting sequences).

Let, $u = 012102010212$ and choose

$$\beta_1 = 00000000000111111101111$$

$$\beta_2 = 000000110100100111101111$$

Then $\beta_1[u] = 012102010210121020120212 = \beta_2[u]$.

Example. A shuffled word $w \in u \sqcup u$ can be square-free even if u is not.

Let $u = (012)^2$ then the following words in $u \sqcup u$ are square-free:

$\beta_0[u] = 012010212012$, where $\beta_0 = 000001011111$

$\beta_1[u] = 010201210212$, where $\beta_1 = 001001101011$

$\beta_2[u] = 010210120212$, where $\beta_2 = 001010011011$

Main result

Theorem. *There exist infinite square-free $u \in \{0,1,2\}^{\mathbb{N}}$ such that $u \sqcup u$ has a square-free word.*

Sketch. The words to be shuffled will be images of the 12-uniform morphism $\rho: \{0,1,2,3\}^* \rightarrow \{0,1,2\}^*$:

$$\rho(0) = 010210120212$$

$$\rho(1) = 012101202102$$

$$\rho(2) = 012102010212$$

$$\rho(3) = 012102120102.$$

$$\rho(0) = 010210120212$$

$$\rho(1) = 012101202102$$

$$\rho(2) = 012102010212$$

$$\rho(3) = 012102120102.$$

- The images $\rho(i)$ are square-free, but
the morphism ρ is not square-free: the images of
12, 20 and 30 have squares.

For instance, $\rho(20) = 01210 \cdot (201021)^2 \cdot 0120212$.

- We need an extra morphism α to fix this problem.

Nevertheless

Each $\rho(i)$ can be shuffled to obtain a square-free word:
 $\sigma(i) \in \rho(i) \sqcup \rho(i)$.

$$\sigma(0) = 010210120102120210120212$$

$$\sigma(1) = 012101202101210201202102$$

$$\sigma(2) = 012102010210121020120212$$

$$\sigma(3) = 012102120102101202120102$$

The fix: α

$$\alpha(0) = 1013$$

$$\alpha(1) = 1023$$

$$\alpha(2) = 1032$$

- For all $w \in \{0,1,2\}^*$, the images $\alpha(w)$ avoids the ‘forbidden’ words 12, 20 or 30.
- The morphism α is surely square-free.

Combine:

$$B(i) = \rho\alpha(i)$$

$$S(i) = \sigma\alpha(i)$$

The images of the words $B(i)$ are:

$$B(0) = 012101202102010210120212012101202102012102120102$$

$$B(1) = 012101202102010210120212012102010212012102120102$$

$$B(2) = 012101202102010210120212012102120102012102010212$$

The lengths of these words are 48.

The images of the shuffled words are of length 96:

$$S(0) = 012101202101210201202102010210120102120210120212 \\ \cdot 012101202101210201202102012102120102101202120102$$

$$S(1) = 012101202101210201202102010210120102120210120212 \\ \cdot 012102010210121020120212012102120102101202120102$$

$$S(2) = 012101202101210201202102010210120102120210120212 \\ \cdot 012102120102101202120102012102010210121020120212$$

By Crochemore's criterion and a computer check:
the morphisms B and S are square-free.

Finally:

$$S(w) \in B(w) \sqcup B(w)$$

for infinite square-free words w . (Details omitted.) \square

Corollary. *There exist infinite square-free ternary words that are Abelian periodic:*

$$w_1 w_2 \cdots \quad \text{where } w_i \sim w_j$$

(same Parikh vectors).

A simpler solution to the corollary was given earlier using stem factorizations.

Open questions

The main question involves self-shuffling.

Problem 1. *Does there exist an infinite square-free word w such that $w \in w \sqcup w$?*

Problem 2. *Which words $w \in u \sqcup u$ can be obtained in more than one way from a single u ?*

Example. $u \sqcup u$ can contain several square-free words.

E.g., $u = 01021201$ gives rise to 3 square-free words, but $v = 01201021$ only one.

There are also square-free words that do not shuffle to any square-free word.

Example. Square-free words in $u \sqcup u$ seem to be relatively sparse compared to the number of all square-free words. **How sparse?**

(Note that the letters have even parity in $u \sqcup u$.)

Length L	# of all square-free	# shuffles $\beta[u]$	# $ u = L/2$
4	18	0	0
6	42	6	6
8	78	12	6
10	144	30	12
12	264	24	18
14	456	42	30
16	798	78	42
18	1392	138	36
20	2388	228	54
22	4146	396	138
24	7032	588	168
26	11892	1008	234